

Resolving sets in finite projective planes

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Joint work with Tamás Héger

Finite Geometry Conference and Workshop

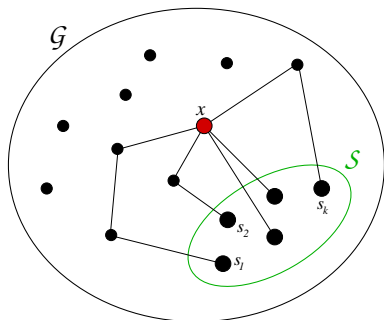
10-14 June 2013

Szeged, Hungary

Let $G = (V, E)$ be a simple graph.

$d(x, y)$: distance of x and y , $x, y \in V$

$S = \{s_1, s_2, \dots, s_k\}$ vertex set



$$d(x, s_1)$$

$$d(x, s_2)$$

...

$$d(x, s_k)$$

x is **resolved** by S if its distance list is different from all the other distance lists

Definition (Resolving set)

The subset $S = \{s_1, \dots, s_k\} \subset V$ is a **resolving set**, if the ordered distance lists $(d(x, s_1), \dots, d(x, s_k))$ are different for all $x \in V$.

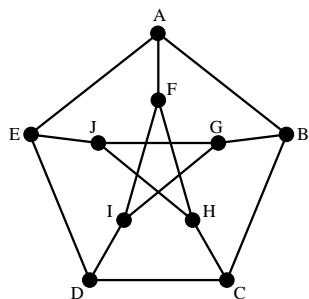
In other words:

$S = \{s_1, \dots, s_k\} \subset V$ is a **resolving set** \iff
 $\forall x, y \in V \exists z \in S: d(x, z) \neq d(y, z)$.

Definition (Metric dimension)

The **metric dimension** $\mu(G)$ is the size of the smallest resolving set.

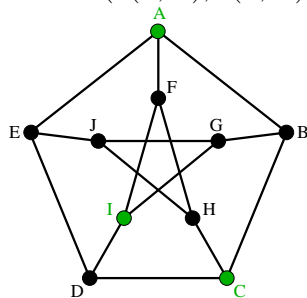
Example: Petersen graph



Example: Petersen graph

Resolving set: A, C, I

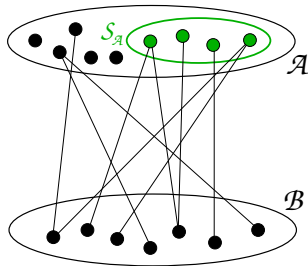
$x \in V: (d(x, A), d(x, C), d(x, I))$



$A :$	$(0, 2, 2)$	$F :$	$(1, 2, 1)$
$B :$	$(1, 1, 2)$	$G :$	$(2, 2, 1)$
$C :$	$(2, 0, 2)$	$H :$	$(2, 1, 2)$
$D :$	$(2, 1, 1)$	$I :$	$(2, 2, 0)$
$E :$	$(1, 2, 2)$	$J :$	$(2, 2, 2)$

Metric dimension: $\mu(\text{Petersen}) = 3$

Let $G = (A \cup B, E)$ be a bipartite graph.
split resolving set:

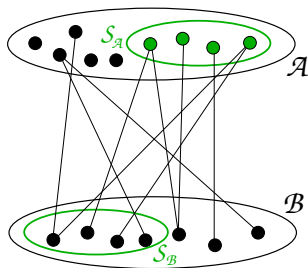


Let $G = (A \cup B, E)$ be a bipartite graph.

split resolving set:

$S_A \subset A$ resolves B

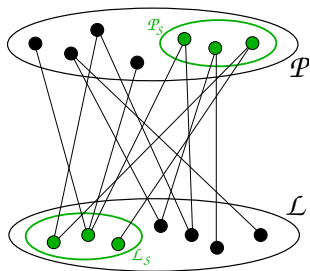
$S_B \subset B$ resolves A



S_A and S_B are called **semi-resolving sets**.

Next talk: Tamás Héger - semi-resolving sets!

Resolving sets in incidence graphs of finite projective planes



(P, ℓ) edge $\Leftrightarrow P \in \ell$

$G = (\mathcal{P}, \mathcal{L}, E)$: incidence graph of a finite projective plane

$S = (\mathcal{P}_S \cup \mathcal{L}_S)$ resolving set in $\Pi_q \iff$

S is a resolving set in the incidence graph

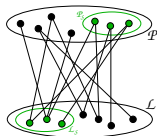
$d(P_1, P_2) = 2$, $d(\ell_1, \ell_2) = 2$, $d(P, \ell) = 1$ or 3

metric dimension $\mu(G)$:

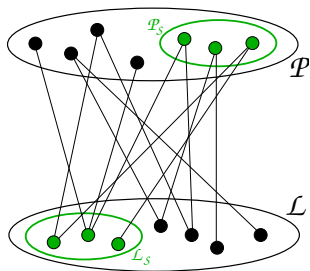
- first introduced by Harary and Melter and (independently) by Slater in the 1970s
- a survey of Bailey and Cameron (2011)
- connection with other graph parameters
- talk of **Robert Bailey** on the 23rd **BCC**, Exeter, July 2011:
- he asked the metric dimension of $G = (\mathcal{P}, \mathcal{L}, E)$
- construction of a resolving set of size $4q - 1$ (Bill Martin)

- connection with other graph parameters: **dimension jump**
- $B \subset V(G)$ is a **base**: the only automorphism of G that fixes B pointwise is the **identity**
- $b(G)$ **base size**: the size of the smallest base of G
- $b(G) \leq \mu(G)$
- **dimension jump**: $\delta(G) = \mu(G) - b(G)$
- projective planes: highly symmetric graphs with large dimension jump
- Γ : incidence graph of $\text{PG}(2, q)$,
order of Γ : $n = 2(q^2 + q + 1)$
- $b(\Gamma) \leq 5$
- $\delta(\Gamma) \geq 4q - 9 \sim 2\sqrt{2n}$

- Π_q : projective plane of order q
- $S = \mathcal{P}_S \cup \mathcal{L}_S$: resolving set in the incidence graph of Π_q
- PQ : line joining two distinct points P and Q
- $[P]$: set of lines through a point P
- $[\ell]$: set of points on a line ℓ
- tangent line \longleftrightarrow 1-covered point
- skew line \longleftrightarrow not covered point



Resolving sets in incidence graphs of finite projective planes



(P, ℓ) edge $\Leftrightarrow P \in \ell$

$G = (\mathcal{P}, \mathcal{L}, E)$: incidence graph of a finite projective plane

$S = (\mathcal{P}_S \cup \mathcal{L}_S)$ resolving set in $\Pi_q \iff$

$\forall \ell \in \mathcal{L}: \ell \in \mathcal{L}_S, \text{ or } \ell \cap \mathcal{P}_S \text{ is unique}$

$\forall P \in \mathcal{P}: P \in \mathcal{P}_S, \text{ or } [P] \cap \mathcal{L}_S \text{ is unique}$

Lemma

Let $S = \mathcal{P}_S \cup \mathcal{L}_S$, ℓ be a line in Π_q . If $|\ell \cap \mathcal{P}_S| \geq 2$ then ℓ is resolved by S .

Dually, let P be a point in Π_q . If $|[P] \cap \mathcal{L}_S| \geq 2$ then P is resolved by S .

- Points and lines in S are resolved
- (At least 2-)secants are resolved
- (At least) 2-covered points are resolved

We have to distinguish:

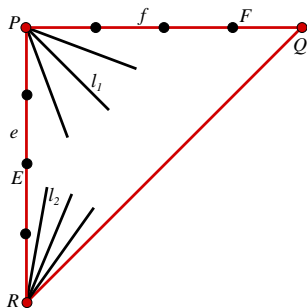
- tangents and skew lines (to \mathcal{P}_S)
- 1-covered points and not covered points (by \mathcal{L}_S)

Proposition

$S = \mathcal{P}_S \cup \mathcal{L}_S$ is a resolving set in a finite projective plane if and only if the following properties hold for S :

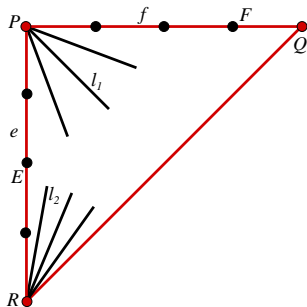
- **P1** There is at most one outer line skew to \mathcal{P}_S .
- **P1'** There is at most one outer point not covered by \mathcal{L}_S .
- **P2** Through every inner point there is at most one outer line tangent to \mathcal{P}_S .
- **P2'** On every inner line there is at most one outer point that is 1-covered by \mathcal{L}_S .

Example



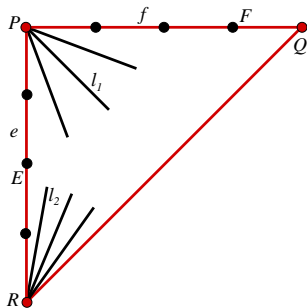
- P1 Outer line skew to \mathcal{P}_S : RQ .
- P1' Outer point not covered by \mathcal{L}_S : Q .
- P2 Through $F \in (f \cap \mathcal{P}_S)$ there is no outer tangent line.
Through $E \in (e \cap \mathcal{P}_S)$ outer tangent line: EQ .
- P2' On $l_1 \in ([P] \cap \mathcal{L}_S)$ outer 1-covered point: $l_1 \cap RQ$
On $l_2 \in ([R] \cap \mathcal{L}_S)$ there is no outer 1-covered point.

Example



- **P1** Outer line skew to \mathcal{P}_S : RQ .
- **P1'** Outer point not covered by \mathcal{L}_S : Q .
- **P2** Through $F \in (f \cap \mathcal{P}_S)$ there is no outer tangent line. Through $E \in (e \cap \mathcal{P}_S)$ outer tangent line: EQ .
- **P2'** On $l_1 \in ([P] \cap \mathcal{L}_S)$ outer 1-covered point: $l_1 \cap RQ$
On $l_2 \in ([R] \cap \mathcal{L}_S)$ there is no outer 1-covered point.

Example



$$|S| = |\mathcal{P}_S| + |\mathcal{L}_S| = 2q - 2 + 2q - 2 = 4q - 4$$

Proposition (T. Héger, M. Takáts)

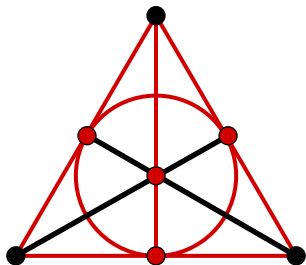
The metric dimension of a projective plane of order $q \geq 23$ is $4q - 4$.

- Classification of the resolving sets of size $4q - 4$, if $q \geq 23$.

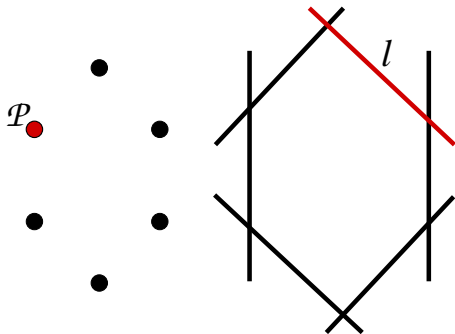
If q is small ($q < 23$):

Fano plane:

$$\mu(\text{PG}(2, 2)) = 5$$



$\mu(\text{PG}(2, 4)) = 10$
construction: hyperoval



Sketch of the proof

- Suppose $S = \mathcal{P}_S \cup \mathcal{L}_S$ is a resolving set, $|S| \leq 4q - 4$.
- Show that $\mu(\Pi_q) = 4q - 4$.
- Characterize the resolving sets of size $4q - 4$.
- combinatorial methods
- we use **duality**

Sketch of the proof

Proposition

$$2q - 5 \leq |\mathcal{P}_S| \leq 2q + 1, 2q - 5 \leq |\mathcal{L}_S| \leq 2q + 1.$$

Proof

t: the number of tangents not in \mathcal{L}_S

Bounds: $t \leq |\mathcal{P}_S|$ and $|\mathcal{P}_S| + |\mathcal{L}_S| \leq 4q - 4$

Double counting the pairs:

$$\star = \{(P, \ell) : P \in \mathcal{P}_S, P \in [\ell], |[\ell] \cap \mathcal{P}_S| \geq 2\}$$

$$2(q^2 + q + 1 - 1 - t - |\mathcal{L}_S|) \leq \star \leq |\mathcal{P}_S|(q + 1) - t$$

We get $2q - 5 \leq |\mathcal{P}_S|$ and dually $2q - 5 \leq |\mathcal{L}_S|$. □

Sketch of the proof

Proposition

Let $q \geq 89$. Then any line intersects \mathcal{P}_S in either ≤ 7 or $\geq q - 4$ points.

Proof

ℓ : secant of S , $|\ell \cap \mathcal{P}_S| = x$, $2 \leq x \leq q$

Let $P \in \ell \setminus \mathcal{P}_S$.

$s(P), t(P)$: number of skew/tangent lines to \mathcal{P}_S through P , resp.

s : the number of skew lines

t : the number of tangents intersecting $\ell \setminus \mathcal{P}_S$

Bounds: $s \leq |\mathcal{L}_S| + 1$ and $s + t \leq |\mathcal{L}_S| + (1 + |\mathcal{P}_S| - x)$

$\implies 2s + t \leq 2|\mathcal{L}_S| + |\mathcal{P}_S| - x + 2;$

$|\mathcal{P}_S| + |\mathcal{L}_S| \leq 4q - 4$

Proof (cont'd)

Counting the points of \mathcal{P}_S on ℓ and on $[P]$:

$$2q + 1 \geq |\mathcal{P}_S| \geq x + t(P) + 2(q - t(P) - s(P))$$

$$x \leq 2s(P) + t(P) + 1$$

Adding up the inequalities $\forall P \in [\ell] \setminus \mathcal{P}_S$:

$$(q + 1 - x)x \leq 2s + t + (q + 1 - x).$$

We get $x^2 - (q + 3)x + 7q \geq 0 \Rightarrow$

$x \leq 7$ or $x \geq q - 4$, if $q \geq 89$.



Sketch of the proof

Proposition

Let $q \geq 47$. Then there exist two lines intersecting \mathcal{P}_S in at least $q - 4$ points.

Proof

Suppose to the contrary: there is ≤ 1 line intersecting \mathcal{P}_S in at least $q - 4$ points.

ℓ : the longest secant, let $x = |[\ell] \cap \mathcal{P}_S| \geq 2$.

Note that $x \leq 7$ is also possible.

n_i : the number of i -secants different from ℓ

Let $n_0 = s$, $n_1 = t$, and let $b = |\mathcal{P}_S|$

The standard equations yield:

$$\sum_{i=2}^7 n_i = q^2 + q + 1 - s - t - 1$$

$$\sum_{i=2}^7 i n_i = (q + 1)b - t - x$$

$$\sum_{i=2}^7 i(i - 1)n_i = b(b - 1) - x(x - 1)$$

Proof (cont'd)

$$\begin{aligned}0 &\leq \sum_{i=2}^7 (i-2)(7-i)n_i = \\&= -\sum_{i=2}^7 i(i-1)n_i + 8\sum_{i=2}^7 in_i - 14\sum_{i=2}^7 n_i = \\&= -b^2 + (8q+9)b + x(x-9) + 6(s+t) + 8s - 14(q^2+q).\end{aligned}$$

Bounds: $s+t \leq |\mathcal{P}_S| + |\mathcal{L}_S| + 1 \leq 4q-3$, $s \leq |\mathcal{L}_S| + 1 \leq 2q+2$
and $x \leq q+1$

we get $0 \leq -b^2 + (8q+9)b - 13q^2 + 19q - 10$.

$$b = |\mathcal{P}_S| \leq 2q+1$$

For $b = 2q+2$:

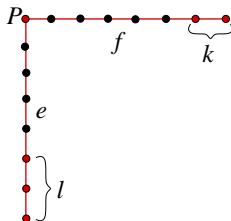
$$-q^2 + 45q - 20 < 0 \text{ if } q \geq 47.$$

Hence $b > 2q+2$, a contradiction. □

Sketch of the proof

There exist two distinct lines $e, f: e \cap f = P$

$|[e] \cap (\mathcal{P}_S \setminus P)| = q - l$ and $|[f] \cap (\mathcal{P}_S \setminus P)| = q - k$ with $k \leq l \leq 5$.



Proposition

Suppose $q \geq 23$. Then $k + l \leq 3$. Moreover, $l = 3$ is not possible.

Dually:

Proposition

There exist two points P, R such that $|[P] \cap \mathcal{L}_S|$ and $|[R] \cap \mathcal{L}_S|$ is at least $q - 4$.

For the points P, R :

$|[P] \cap \mathcal{L}_S| = q - l'$ and $|[R] \cap \mathcal{L}_S| = q - k'$ with $k', l' \leq 5$.

Proposition

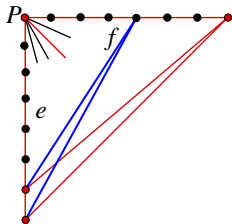
$k' + l' \leq 3$. Moreover, $k' = 3$ or $l' = 3$ is not possible.

Suppose: $|\mathcal{P}_S| \leq |\mathcal{L}_S|$.

$\implies |\mathcal{P}_S| \leq 2q - 2$.

Sketch of the proof

- $k + l \leq 3 \implies |\mathcal{P}_S| \geq 2q - 3$
- $|S| = |\mathcal{P}_S| + |\mathcal{L}_S| \leq 4q - 4$, we assumed $|\mathcal{P}_S| \leq |\mathcal{L}_S|$
- If $|\mathcal{P}_S| = 2q - 3$ then $|\mathcal{L}_S| \geq 2q - 1$

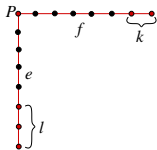


- If $|\mathcal{P}_S| = 2q - 2$ then $|\mathcal{L}_S| \geq 2q - 2$

Corollary

Suppose $q \geq 23$. The metric dimension of Π_q is $\mu = 4q - 4$.

Sketch of the proof



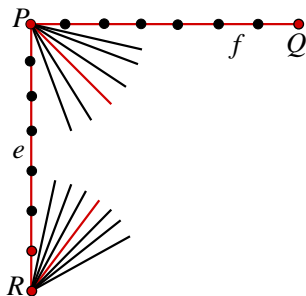
Proposition

*There exists a point $R \in e \setminus P$ such that $|([R] \setminus e) \cap \mathcal{L}_S| \geq q - 1$.
Moreover, if $l = 2$ then $R \notin \mathcal{P}_S$.*

Proposition

*Let e, f be the lines such that $e \cap f = P$, $|[e] \cap (\mathcal{P}_S \setminus P)| = q - l$
and $|[f] \cap (\mathcal{P}_S \setminus P)| = q - k$, $k + l \leq 3$, $l \neq 3$. Then
 $|([P] \setminus \{e, f\}) \cap \mathcal{L}_S| \geq q - 2$.*

So we get the following:

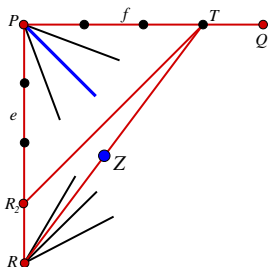
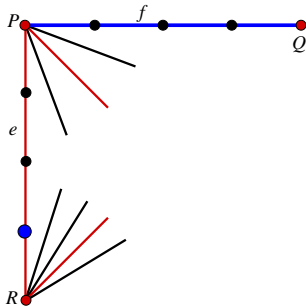
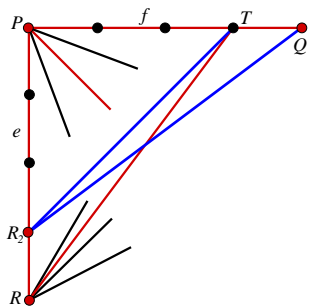


And we need 2 more objects in addition:

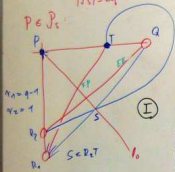
2 lines or 1 point and 1 line

surprisingly many (more than 30) different types

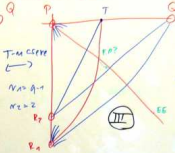
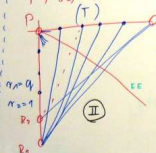
Some constructions



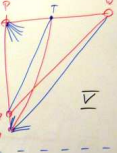
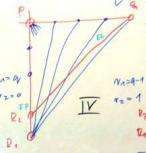
$P \in \mathcal{P}_s$
 $|\mathcal{P}_s| = 2q - 2$



$P \notin \mathcal{P}_s$, $k+l=3$ & $|\mathcal{P}_s| = 2q - 3$

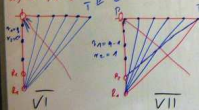


T.G.

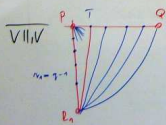
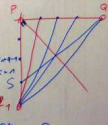
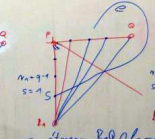
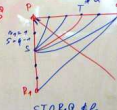
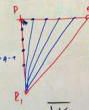
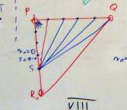


$P \notin \mathcal{P}_s$, $|\mathcal{P}_s| = 2q - 2$, $k+l=2$

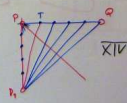
a) $k=0, l=2$



b) $k=1, l=1$

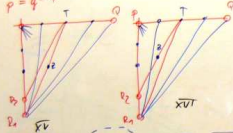


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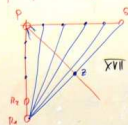


$p \notin P_5, |P_5| = 2q-2, q+l=3$ von links part $z \in P_5$ ($v_1 \geq q-1$)

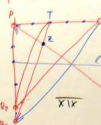
$p = q-1$



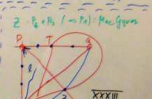
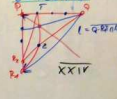
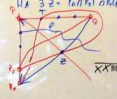
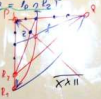
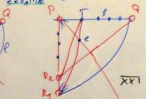
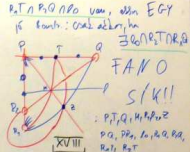
$p = q-2, v_1 = q$



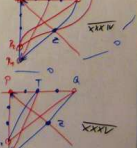
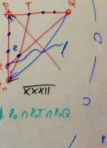
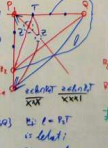
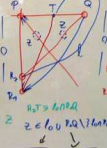
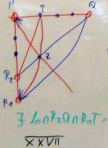
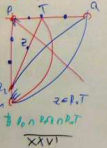
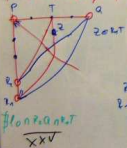
$p = q-2, v_1 = q-1, v_2 = 0$



$z \in P_5, T \in P_5$



$p = q-2, v_1 = q-1, v_2 = 1$



Thanks for your attention!